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Code Injection via Arbitrary Pointer Overwrite

Introduction – Eduard Muresan

- **L** Security Researcher @ Bitdefender
- L Passionate about low level systems programming
- Ly Interested in injection techniques
- L Mythic raider at night

Agenda

- L Process Injection Background
- L C.I.A.P.O. Methodology
- L Executable Pointer Examples
- L Mitigations
- L Demo
- L, Q&A

What is process injection

Process injection

A method used by malware to execute arbitrary code within the address space of a separate live process.

Why process injection

Process injection

Execution via process injection might evade process-based detections from security products.

Injection building blocks

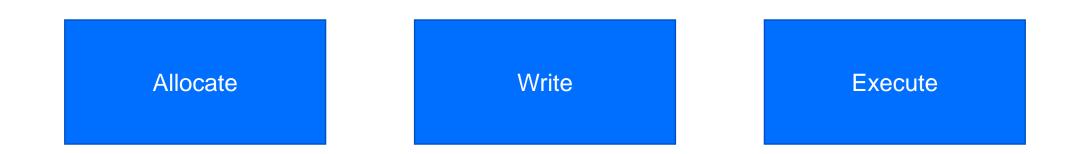
Allocate

Injection building blocks

Allocate

Write

Injection building blocks



Injection building blocks

Allocate

Usually legitimateCan be implicit

Write

Execute

Injection building blocks

Allocate

Write

Might be legitimateCan be implicit

Execute

Injection building blocks

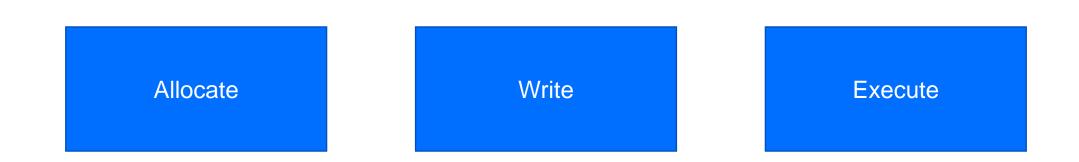


Write

Execute

Not so legitimateHighly monitored

Piecing everything together



Piecing everything together

```
void RemoteThreadInjection(HANDLE Process) {
    // Pretend this is a fancy shellcode.
    static const BYTE shellcode[] = { 0xC3 };
```

// Reserve space for the payload.
void* payload = VirtualAllocEx(ProceAllocateLL, sizeof(shellcode), MEM_RESERVE |
 MEM_COMMIT, PAGE_EXECUTE_READWRITE);

// Write the payload to the allocated space.
Write
Write
WriteProcessMemory(Process, payload, shellcode, sizeof(shellcode), NULL);

// Trigger the payload.
CreateRemoteThread(Process, NULL, 0, payload, NULL, 0, NULL);

Piecing everything together

```
void RemoteThreadInjection(HANDLE Process) {
    // Pretend this is a fancy shellcode.
    static const BYTE shellcode[] = { 0xC3 };
```

```
// Reserve space for the payload.
void* payload = VirtualAllocEx(Process, NULL, sizeof(shellcode), MEM_RESERVE |
    MEM_COMMIT, PAGE_EXECUTE_READWRITE);
```

// Write the payload to the allocated space.
WriteProcessMemory(Process, payload, shellcode, sizeof(shellcode), NULL);

```
// Trigger the payload.
CreateRemoteThread(Process, NULL, 0, payload, NULL, 0, NULL);
```

Change the execute primitive

```
void APCInjection(HANDLE Process, HANDLE Thread) {
    // Pretend this is a fancy shellcode.
    static const BYTE shellcode[] = { 0xC3 };
    // Reserve space for the payload.
    void* payload = VirtualAllocEx(Process, NULL, sizeof(shellcode), MEM_RESERVE | MEM_COMMIT,
        PAGE EXECUTE READWRITE);
```

// Write the payload to the allocated space.
WriteProcessMemory(Process, payload, shellcode, sizeof(shellcode), NULL);

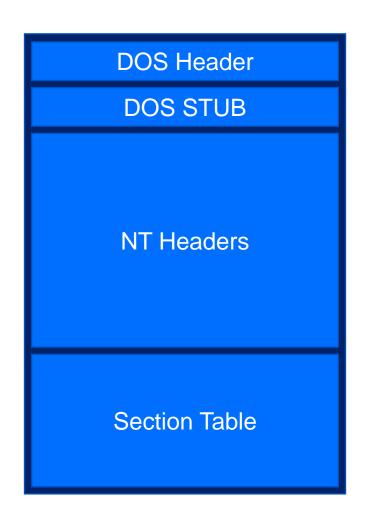
```
// Trigger the payload.
QueueUserAPC(payload, Thread, NULL);
```

The problem with execute primitives

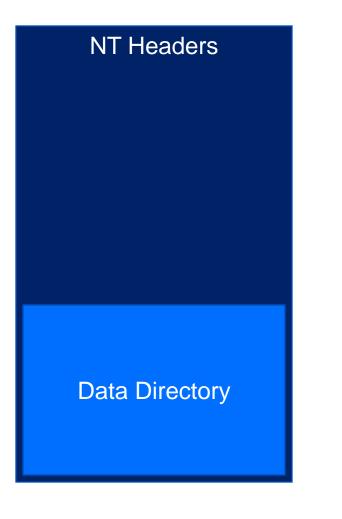
The problem with execute primitives

Highly monitored
Can be blocked
Might be traced back to the attacker
Becomes a cat and mouse game

Remove the execute primitive?



NT Headers	



Thread Local Storage Injection



```
typedef struct _IMAGE_TLS_DIRECTORY64 {
   ULONGLONG StartAddressOfRawData:
   ULONGLONG EndAddressOfRawData;
   ULONGLONG AddressOfIndex; // PDWORD
   ULONGLONG AddressOfCallBacks; // PIMAGE_TLS_CALLBACK *;
   DWORD SizeOfZeroFill;
   union {
       DWORD Characteristics;
       struct {
           DWORD Reserved0 : 20;
           DWORD Alignment : 4;
           DWORD Reserved1 : 8;
        } DUMMYSTRUCTNAME;
    } DUMMYUNIONNAME;
} IMAGE_TLS_DIRECTORY64;
```

NT Headers	



Entry Point Injection

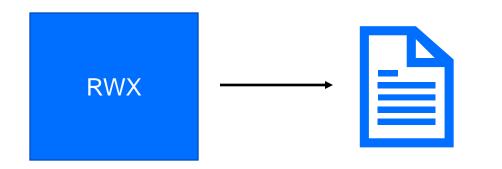
NT Headers

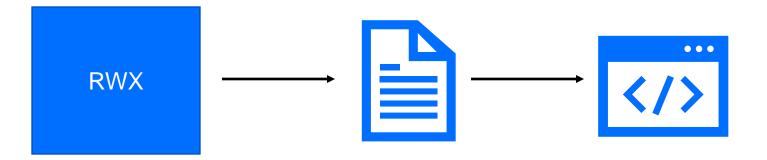
AddressOfEntryPoint

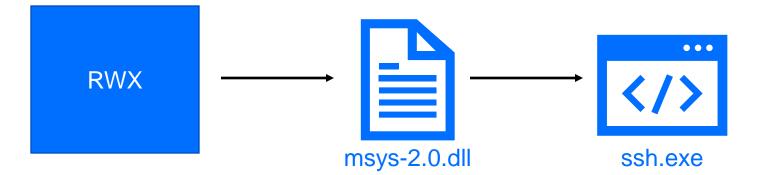
```
typedef struct _PEB_LDR_DATA {
    BYTE Reserved1[8];
    PVOID Reserved2[3];
    LIST_ENTRY InMemoryOrderModuleList;
} PEB_LDR_DATA, * PPEB_LDR_DATA;
```

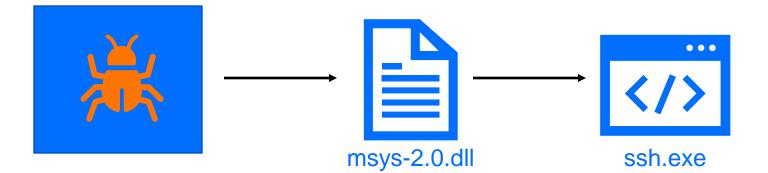
typedef struct _LDR_DATA_TABLE_ENTRY {
 PVOID Reserved1[2];
 LIST_ENTRY InMemoryOrderLinks;
 PVOID Reserved2[2];
 PVOID DllBase;
 PVOID EntryPoint;
 PVOID Reserved3;
 UNICODE_STRING FullDllName;
 [snip]
} LDR_DATA_TABLE_ENTRY, * PLDR_DATA_TABLE_ENTRY;







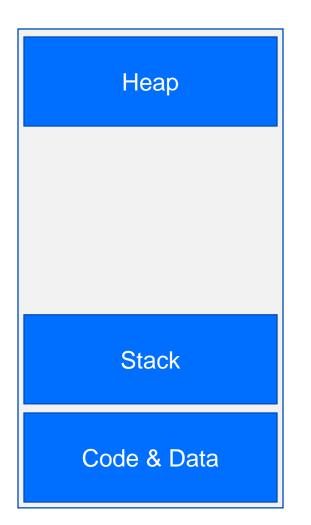


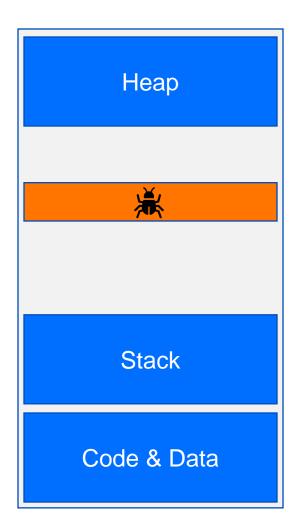


The problem with specific pointers

The problem with specific pointers

Highly monitoredExecution can be blockedBecomes a cat and mouse game

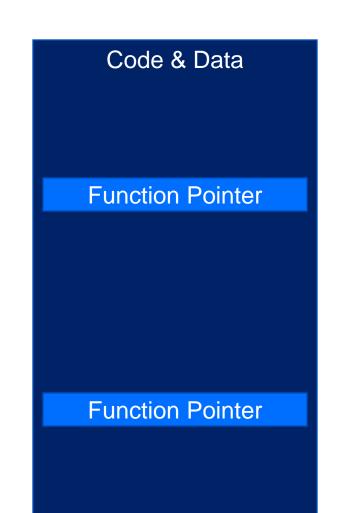






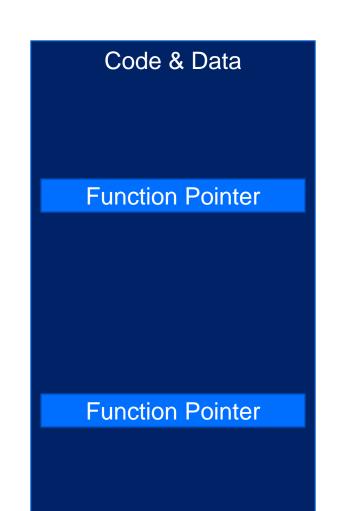
C.I.A.P.O. Methodology

Heap **Function Pointer Function Pointer Function Pointer**



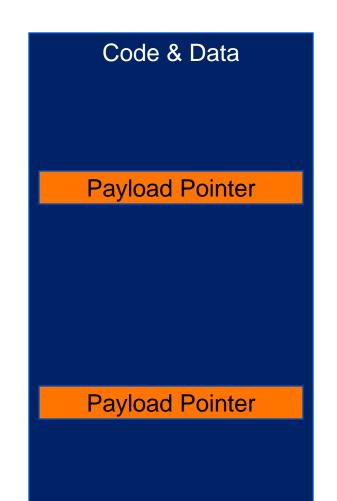
C.I.A.P.O. Methodology

Heap **Payload Pointer Function Pointer Function Pointer**



C.I.A.P.O. Methodology

Heap **Payload Pointer Payload Pointer Payload Pointer**



All your pointer are belong to us



Example

```
static void (*pSleep)(DWORD);
int main() {
    pSleep = GetProcAddress(
        GetModuleHandleA("kernel32.dll"),
        "Sleep"
    );
    while (TRUE) {
        pSleep(1000);
    }
    return 0;
```

Example

```
static void (*pSleep)(DWORD);
int main() {
    pSleep = GetProcAddress(
        GetModuleHandleA("kernel32.dll"),
        "Sleep"
    );
    while (TRUE) {
        pSleep(1000);
    }
    return 0;
```

sub	rsp, 28h
lea	<pre>rcx, ModuleName ; "kernel32.dll"</pre>
call	<pre>cs:imp_GetModuleHandleA</pre>
mov	rcx, rax ; hModule
lea	rdx, ProcName ; "Sleep"
call	<pre>cs:imp_GetProcAddress</pre>
mov	cs:pSleep, rax
loc_140 mov	001028: ecx, 3E8h
call	rax
mov	rax, cs:pSleep
jmp	short loc_140001028

Indirect Branches

- Dynamically determine the target address at runtime
- Generated by the compiler in most programs
- Almost guaranteed to be executed at some point

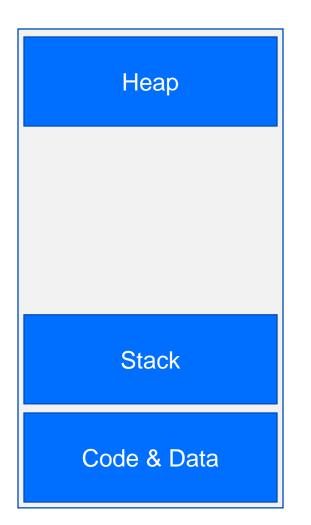
sub	rsp, 28h
lea	<pre>rcx, ModuleName ; "kernel32.dll"</pre>
call	cs:imp_GetModuleHandleA
mov	rcx, rax ; hModule
lea	rdx, ProcName ; "Sleep"
call	<pre>cs:imp_GetProcAddress</pre>
mov	cs:pSleep, rax
loc_140	001028:
mov	ecx, 3E8h
call	rax
mov	rax, cs:pSleep
jmp	short loc_140001028

Example

Autos	→ 및 (×
Search (Ctrl+E)	ho - $ ightarrow$ Search Depth: 3 -	
Name	Value Type	4
🗊 pSleep	0x00007ffdbf59b0b0 {kernel32.dll!SleepStub} void(*)(unsigned long)	
Autos Locals	Watch 1	

Example

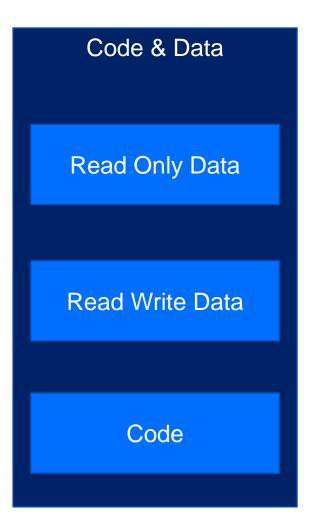
Autos				- 4	х
Search (Ctrl+E)	- م _ر	$\leftrightarrow ightarrow$ Search Dept	th: 3 •	•	
Name	Value		Туре		-
🗊 pSleep	0x000001b27cf90080		void(*)(ur	nsigned long)	
Autos Locals	Watch 1				







Finding executable pointers



Read Only Data:
Import Address Table
Virtual Function Tables
Overwriting might be suspicious

Finding executable pointers

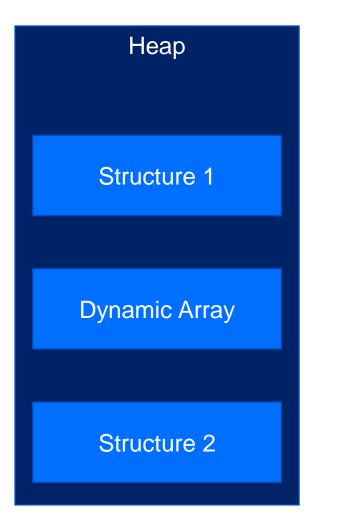


Read Write Data: Arbitrary function pointers

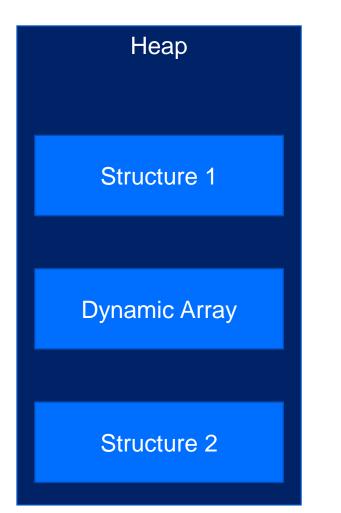


_GetPr	ocFromComCtl32
mov	[rsp+arg_0], rbx
mov	[rsp+arg_8], rsi
push	rdi
sub	rsp, 20h
mov	rdi, rcx
xor	ebx, ebx
mov	<pre>rcx, cs:g_hinstCC</pre>
mov	rsi, rdx
	rcx, rcx
jz	short loc_14008D2B1
loc_14	008D28B:
mov	rdx, rsi
call	<pre>cs:imp_GetProcAddress</pre>
nop	dword ptr [rax+rax+00h]
mov	rbx, rax
loc_14	008D29D:
mov	[rdi], rbx
mov	rbx, [rsp+28h+arg_0]
mov	rsi, [rsp+28h+arg_8]
add	rsp, 20h
рор	rdi
retn	



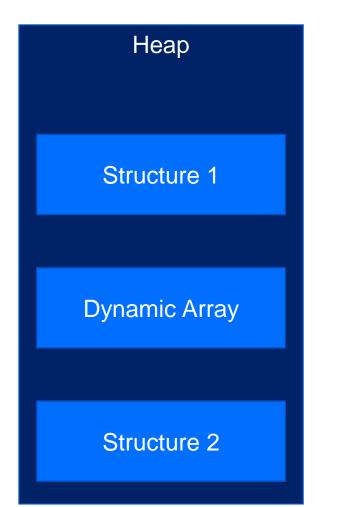


Finding executable pointers



The heap: Arbitrary function pointers Usually contained in various structures

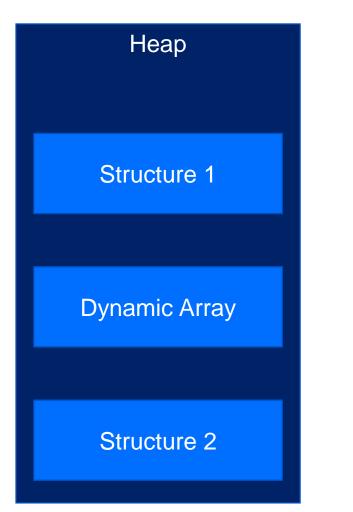
Finding executable pointers



TppAlpcpExecuteCallbackTppTimerQueueExpirationLdrpWorkCallback

Thread Pool Objects:
All processes have a thread pool
Represented as structures
Multiple work item types
Timers, Workers and async Workers (Wait, IO, Alpc)

Finding executable pointers



Kernel32!SortCompareString National Language Support Local specific string compare Loaded based on registry key

Mitigations: CET-IBT

JMP RAX

Mitigations: CET-IBT



Mitigations: CET-IBT



Mitigations: CFG

while (TRUE) {
 pSleep(1000);
 }

Mitigations: CFG

while (TRUE) { pSleep(1000); }

loc_14	0001028:
mov	ecx, 3E8h
call	rax
mov	rax, cs:pSleep
jmp	short loc_140001028

Mitigations: CFG

while (TRUE) { pSleep(1000);

loc_140	0001030:
mov	ecx, 3E8h
call	<pre>cs:guard_dispatch_icall_fptr</pre>
mov	rax, cs:pSleep
jmp	short loc_140001030

Demo

C 💷 🎟 🗄 🍕 🗙 🔎	ا ی				MA		<filter by="" name<="" th=""><th>C:\Users\emuresan\Desktop></th><th></th></filter>	C:\Users\emuresan\Desktop>	
ess	PID	CPU	Private Bytes	Working Set Cont		1.	114.5		
Registry	124		11,120 K	95,464 K n/a					
System Idle Process	0	98.51	56 K	8 K n/a					
System	4	< 0.01	192 K	148 K n/a					
csrss.exe	508		1,944 K	5,540 K n/a					
wininit.exe	584		1,616 K	7.284 K n/a					
csrss.exe	592	< 0.01	2,340 K	5,724 K n/a					
winlogon.exe	680		2,952 K	12,952 K n/a					
fontdrvhost.exe	896		3,228 K	7,208 K n/a					
dwm.exe	820		52,444 K	87,916 K n/a					
explorer.exe	4920	0.49	56,100 K	130,484 K CFG					
SecurityHealthSystray.exe	7512		1,988 K	9,440 K CFG					
OneDrive.exe	7620		48,244 K	110,080 K CFG					
C msedge.exe	8084		50,916 K	120,816 K CFG				×	
e msedge.exe	8128		2,108 K	7,660 K CFG					
e msedge.exe	6020 6064		10,768 K	25,260 K CFG					
emsedge.exe	6064		11,384 K 7,368 K	31,732 K CFG 18,084 K CFG					
msedge.exe	8636		59,924 K	94,956 K CFG					
msedge.exe	8832		21,200 K	27,988 K CFG					
procexp64.exe	3696	0.49	24,848 K	46,104 K					
cmd.exe	5940	0.45	4,392 K	4,216 K CFG					
conhost.exe	6988	< 0.01	10,384 K	25,636 K CFG					
vpnui.exe	9024		17,048 K	36,060 K CFG					
Usage: 1.47% Commit Charge:	10.75%	2							

Takeaways

- Detecting injection attacks by only monitoring execute primitives is necessary but not sufficient.
- Process Injection allow attackers to evade security solutions, increasing the Mean Time to Detection.
- Having a comprehensive approach to detecting and preventing injections is mandatory to ensure individuals and organizations remain unharmed via Defense in Depth.
- C.I.A.P.O. pushes forward the state of the art, demonstrating that attackers might use brute force approaches to overwrite random functions.





Trusted. Always.

